Theory and practice of algorithmic self-assembly

Damien Woods



Paris

ER02 - ENS Lyon - Jan 2017 - Lectures 3 & 4

## Structure

- Monday (lecture 1). Some high-level motivations, basic algorithmic selfassembly models (definitions) and very recent results on implementing algorithmic DNA nanotube circuits, a self-assembly model, in the wet-lab
- Tuesday (lecture 2): DNA sequence design and results on the DNA nanotube circuit model
- Wednesday (lecture 3). Complexity theory for self assembly.
  - Theorem: The (cooperative, temperature >= 2) abstract tile assembly model is intrinsic universal
- Thursday (lecture 4). Complexity theory for self assembly.
  - **Theorem**: The noncooperative (temperature 1) abstract tile assembly model does not simulate the cooperative model

## Square tiles

- finite set of tile types, unlimited supply of each type, non-rotatable
- Each side has a glue (colour) and strength (0,1,2,3,...)
- System has a temperature (e.g. 2)
- Simple local binding rule: A tile sticks to an assembly if enough of its glues match so that the sum of the strengths of the matching glues is at least the temperature

### Tile assembly system: tile set, seed tile, temperature





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 Efficient assembly of scaled complicated connected shapes using a number of tile types roughly equal to the Kolmogorov complexity of the shape Soloveichik, Winfree. SICOMP 2007<sub>6</sub>

### Theory of algorithmic self-assembly

- Helps us understand the abilities and limitations of self-assembly
  - Also, its fun!
- aTAM is Turing universal: can "run" any algorithm Winfree, PhD Thesis 1998.
- Efficiently assemble n x n squares and other simple shapes
  - Using only Θ(log *n*/log log *n*) tile types
- Efficiently assemble arbitrary finite shapes
  - Number of tile types is roughly the Kolmogorov complexity of the shape
    Soloveichik, Winfree. SICOMP 2007
- aTAM is intrinsically universal: there is one tile set that can simulate any tile assembly system
  - Shape complexity can be put into the seed

Doty, Lutz, Patitz, Schweller, Summers, Woods. FOCS 2012

 Thinking about these topics leads to a kind of "complexity theory" for self-assembly



Soloveichik & Winfree, SIAM J Comp, 2007

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### Theory of algorithmic self-assembly

 Helps us understand the abilities and limitations of self-assembly [9] - Also, its fun! 2HAM,  $\tau = c^3$  (IU [9]) aTAM is Turing universal: can "run" any algorithm Winfree. PhD Thesis 1998. Efficiently assemble n x n squares and other [9] しり simple shapes 2HAM,  $\tau = c^2$  (IU [9]) polygon - Using only Θ(log *n*/log log *n*) tile types TAM,  $\tau = 2$ Rothemund, Winfree. STOC 2000 [9] • Efficiently assemble arbitrary finite shapes UII [8] Number of tile types is roughly the Kolmogorov 2HAM,  $\tau = 2$  (IU [9]) *complexity* of the shape hexagon Soloveichik, Winfree. SICOMP 2007 TAM,  $\tau = 2$ [9, 3] aTAM is intrinsically universal: there is one tile set that can simulate any tile assembly system aTAM,  $\tau > 2$  (IU [11]) Shape complexity can be put into the seed Doty, Lutz, Patitz, Schweller, Summers, Woods. FOCS 2012 aTAM,  $\tau =$ aTAM  $\tau = 2$  (IU [12]) Thinking about these topics leads to a kind of Woods, MCU, 2013 "complexity theory" for self-assembly



 Efficient assembly of scaled connected shapes using a number of tile types roughly equal to the Kolmogorov complexity of the shape <u>Soloveichik</u>, Winfree. SICOMP 2007<sub>8</sub>





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- Many other theoretical questions have been asked
- What questions would you ask?
- The goal is to understand the capabilities of these systems!
- Another goal is to motivate what we should build in the lab!
- Next slide: Let's ask a question

### Intrinsic universality

# Is there a set of **intrinsically universal tiles**: a set of aTAM tiles U that can act like **any** other tile set?



One universal tile set to do everything

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### What does "act like" mean?



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Conway's Game of Life is an intrinsically universal cellular automaton



http://otcametapixel.blogspot.com/ http://www.youtube.com/watch?v=xP5-ileKXE8

Durand, Roka, The game of life: universality revisited. 1999

### Comparing tile assembly models

Is there a set of **intrinsically universal tiles** that can **simulate any tile set**?



- What is it that tile assembly systems do?
  - Make shapes and patterns
  - Carry out a crystal-like growth process (dynamics)
- Let define **simulate** using these criteria that are intrinsic to the model
Comparing tile assembly models

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- For (any) simulated tile assembly system  ${\cal T}$ 
  - $T = (\text{tileset T}, \text{ seed assembly } \sigma, \text{ temperature } \tau)$
- Tile assembly system  $\mathcal U$  simulates  $\mathcal T$  if:
  - Tiles from  $\mathcal{T}$  are represented by  $m \ge m \ge m$  supertiles in  $\mathcal{U}$
  - Assemblies produced by U represent exactly assemblies produced by T
    (via a representation function R : Blocks of tiles from U -> tiles from T)
  - **Dynamics are equivalent** in  $\mathcal{U}$  and  $\mathcal{T}$ , ignoring  $m \ge m$  scaling



#### Simulator system



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  - **Dynamics are equivalent** in  $\mathcal{U}$  and  $\mathcal{T}$ , ignoring  $m \ge m$  scaling



- For (any) simulated tile assembly system  ${\cal T}$ 
  - $T = (\text{tileset T}, \text{ seed assembly } \sigma, \text{ temperature } \tau)$
- Tile assembly system  $\mathcal U$  simulates  $\mathcal T$  if:
  - Tiles from  $\mathcal{T}$  are represented by  $m \ge m \ge m$  supertiles in  $\mathcal{U}$
  - Assemblies produced by U represent exactly assemblies produced by T
    (via a representation function R : Blocks of tiles from U -> tiles from T)
  - **Dynamics are equivalent** in  $\mathcal{U}$  and  $\mathcal{T}$ , ignoring  $m \ge m$  scaling



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    (via a representation function R : Blocks of tiles from U -> tiles from T)
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U



- Green tiles are simulated by supertiles
- For each assembly sequence in the simulated tile system, there is an assembly sequence in the simulator, and vice-versa



**Damien Woods** 

Preassembled seed structure (encodes simulated TAS)

U

Preassembled

seed structure

(encodes simulated

TAS)



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• For each assembly sequence in the simulated tile system, there is an assembly sequence in the simulator, and vice-versa



Preassembled

seed structure

(encodes simulated

TAS)



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Preassembled

seed structure

(encodes simulated

TAS)



simulated tile system, there is an assembly

sequence in the simulator, and vice-versa





cIc C e

В

Seed c





**Damien Woods** 

TAS)









Damien Woods



**Damien Woods** 

Α

В

A



**Damien Woods** 

Seed d d D e

A

В

A

a a

cIc C e















**Damien Woods** 



**Damien Woods** 



**Damien Woods** 



**Damien Woods** 



**Damien Woods** 



**Damien Woods** 

Ignoring  $m \ge m$  scaling, production & dynamics are equivalent in the simulated system and simulator









- Green tiles are simulated by supertiles
- For each assembly sequence in the simulated tile system, there is an assembly sequence in the simulator, and vice-versa



**Damien Woods** 

Preassembled seed structure (encodes simulated TAS)

U

Preassembled

seed structure

(encodes simulated

TAS)



- Green tiles are simulated by supertiles
- For each assembly sequence in the simulated tile system, there is an assembly sequence in the simulator, and vice-versa



U

Preassembled

seed structure

(encodes simulated

TAS)



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U

Preassembled

seed structure

(encodes simulated

TAS)





**Damien Woods** 

Seed d d D e e E



**Damien Woods** 





**Damien Woods** 

Seed d d D e e E



**Damien Woods** 



Damien Woods



**Damien Woods** 

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**Damien Woods** 



**Damien Woods** 

15



**Damien Woods** 



**Damien Woods** 



**Damien Woods** 



**Damien Woods** 



**Damien Woods** 



**Damien Woods** 



**Damien Woods** 



**Damien Woods** 

Ignoring  $m \ge m$  scaling, production & dynamics are equivalent in the simulated system and simulator







## Is the abstract tile assembly model intrinsically universal?



# Is the abstract tile assembly model intrinsically universal? Yes!



## **Theorem**: There is a single intrinsically universal tile set *U* that simulates *any* tile assembly system Doty, Lutz, Patitz, Schweller, Summers, Woods. FOCS 2012

**Damien Woods** 

#### Simulation

- For (any) simulated tile assembly system  ${\cal T}$ 
  - $T = (\text{tileset T}, \text{ seed assembly } \sigma, \text{ temperature } \tau)$
- Tile assembly system  $\mathcal{U}$  simulates  $\mathcal{T}$  if:

 $\sim$  Tiles from  ${\mathcal T}$  are represented by  ${m m} \ {m x} \ {m m}$  supertiles in  ${\mathcal U}$ 

- Assemblies produced by  $\mathcal{U}$  represent exactly assemblies produced by  $\mathcal{T}$  (via a representation function R : Blocks of tiles from  $U \rightarrow$  tiles from T)
- Dynamics are equivalent in  $\mathcal{U}$  and  $\mathcal{T}$ , ignoring  $m \ge m$  scaling



#### Superside



frame glu	e tile loc	kup table	blank prob	pe region pro	obe table t	ile lookup table	glue fr	rame
4 O(log	T ) O( T	$4 \log  T $ (	D( T  <sup>2</sup> ) O	$O( T ^2)$	$O( T ^2)$	$O( T ^4 \log  T ) \qquad O($	log   <i>T</i>  )	4

[*T*] is number of tiles in the simulated tileset *T*.



#### Superside















22






#### One-sided binding with a single strength-t south superside



One-sided binding with a single strength-t south superside



One-sided binding with a single strength-t south superside



One-sided binding with a single strength-t south superside



One-sided binding with a single strength-t south superside



## Crawler doing a tile lookup

















24

## A key problem



Better luck next time!



			-					1
	alua	tile lookup	probe	probe region	blook	tile lookup	alua	
1	giue	table	table	probe region	Dialik	table	giue	1
				•				





									1
	glue	tile lookup table	blank	probe region	probe table	tile lookup table	glue		
									i.
								1	1
Damien Woods									
Dufficit Woods									

i i								
	alua	tile lookup	probe	probe region	blook	tile lookup	alua	
	giue	table	table	probe region	Dialik	table	giue	
				- · · · · · · · · · · · · · · · · · · ·				



	glue	tile lookup table	blank	probe region	probe table	tile lookup table	glue	-	
Damien Woods									



Two-sided binding with opposite cooperating supersides



26





2

b

b 3 a

ā 1



2

b

b 3 a

ā 1



2

b

b 3 a

ā 1



2

b

b 3 a

a 1

#### 3-sided "uh-oh" example: probes miss each other











3-sided "uh-oh" example: probes miss each other









3-sided "uh-oh" example: probes miss each other









27

#### 3-sided "uh-oh" example: probes miss each other



3-sided "uh-oh" example: probes miss each other



3 b

?

?

- Variety of cases for different orders of superside arrival
- Superside win/lose configurations and crawler initiation locations (green)
- Proof analogy:
  - Distributed game
  - Computation & geometry
  - Key challenge: make all the tricks work together



**Damien Woods** 





# Is the abstract tile assembly model intrinsically universal?



# Is the abstract tile assembly model intrinsically universal? Yes!



# **Theorem**: There is a single intrinsically universal tile set *U* that simulates *any* tile assembly system Doty, Lutz, Patitz, Schweller, Summers, Woods. FOCS 2012

**Damien Woods** 

## A zoo of self-assembly models

How to compare these models?

abstract tile assembly model



## A zoo of self-assembly models

#### How to compare these models?

#### abstract tile assembly model

temperature 2



locally consistent

temperature 1

. . .

## A zoo of self-assembly models

#### How to compare these models?

hierarchical

abstract tile assembly model

temperature 2



locally consistent

temperature 1

. . .




#### dupled & restricted glue



Intrinsic universality...requires cooperation. Meunier, Patitz, Summers, Theyssier, Winslow, Woods. SODA 2014





## Magic dust



http://nighthawk101stock.deviantart.com/

Result 1: There is a **single** rotatable polygon that simulates **all** tile assembly systems



## Result 2: For each (e.g. Wang) plane tile system there is one rotatable polygon that simulates it

**Theorem 7.1** Each colored square and hexagon plane tiling system in the families  $(\{\}, c_m), (\{t_r, t_f\}, c_m), (\{\}, c_c)$  and  $(\{t_r, t_f\}, c_c)$  is simulated by an n-gon nearly-plane tiling system.



## Theorem: For each (Wang) plane tiling system there is one rotatable polygon that simulates it



Portion of a tiling of Robinson's 10-tile aperiodic square tile set (with rotations)

- Wang plane tiling system:
  - Try to fill the plane with tiles.
  - All sides must match.
  - We care about the existence of tilings, but not how we made the tiling

Demaine, Demaine, Fekete, Patitz, Schweller, Winslow, Woods. One Tile to Rule Them All: Simulating Any Turing Machine, Tile Assembly System, or Tiling System with One Rotatable Puzzle Piece. ICALP 2014.

### A rotatable polygon that simulates a tile set

 For each set of (possibly rotatable, flipable) Wang square/hexagon tiles there is a single rotatable tile that simulates it

Hexagon example

![](_page_188_Figure_3.jpeg)

### A rotatable polygon that simulates a tile set

 For each set of (possibly rotatable, flipable) Wang square/hexagon tiles there is a single rotatable tile that simulates it

![](_page_189_Figure_2.jpeg)

![](_page_189_Figure_3.jpeg)

rotation 1 = t1 rotation 2 = t2

![](_page_189_Figure_5.jpeg)

Robinson's 10-tile aperiodic tile set (complimentary matching constraint)

![](_page_189_Picture_7.jpeg)

![](_page_189_Figure_8.jpeg)

### A rotatable polygon that simulates a tile set

• For each set of (possibly rotatable, and/or flipable) square or hexagon tiles there is a single (rotatable, flipable) tile that simulates it

![](_page_190_Figure_2.jpeg)

- An aperiodic tile set with 1 tile!
- Small gaps (< 1 tile in size) in the tilings
- We have given a general method (a compiler) to convert any square/hexagon plane tiling tile set to a single tile that simulates it

![](_page_191_Figure_3.jpeg)

![](_page_191_Picture_5.jpeg)

### 1 aperiodic tile

- An aperiodic tile set with 1 tile!
- Small gaps in the tilings
- We have given a general method (a compiler) to convert any square/hexagon plane tiling tile set to a single tile that simulates it

### 1 aperiodic tile

- Socolar–Taylor disconnected tile. 2012
- Aperiodic
- Rotations + flips

![](_page_192_Picture_8.jpeg)

![](_page_192_Figure_9.jpeg)

 Open question: Is there a single aperiodic connected 2D tile that makes gap-free tilings of the plane?

### But this is not (yet) magic dust

![](_page_193_Picture_1.jpeg)

### A harder challenge: one tile for all of tile self-assembly

- Simulating tile assembly systems is significantly trickier than plane (Wang) tiling systems
- We want to design a single rotatable, flipable tile that simulates any tile assembly system (Note that as a corollary this gives a single tile that simulates any algorithm)
- Problem 1:
  - Strength tau glues on rotatable tiles => Argh! There's pumpable junk everywhere!
  - Maybe we could find an intrinsically universal square tile set with strength < T glues? No! Any such tile set with a finite seed can not leave the seed's bounding box
  - Lets try hexagons!

### Low strength hexes simulate high-strength squares

Strength 1 or 0 hexagon glues, simulating strength 2, 1 or 0 square glues

![](_page_194_Picture_2.jpeg)

- Then we can simulate a set of low-strength hexagons with a single rotatable polygon
  - Bumps and dents to stop incorrect orientations and incorrect bindings
  - Glues are carefully rearranged on the polygon to allow "self seeding"
  - Many details omitted!

### **Construction overview**

![](_page_195_Figure_1.jpeg)

To use The One, simply apply a sequence of tile assembly system simulations:

![](_page_196_Figure_1.jpeg)

To use The One, simply apply a sequence of tile assembly system simulations:

A tile assembly Tile assembly Hexagonal tile system that simulates Tile assembly system  $U_T$  over the assembly system T using the single system T (with low strength intrinsically rotatable flipable universal tile set U glues) polygonal tile ° C' d C C e Seed

Intrinsically universal tile set

**Damien Woods** 

![](_page_197_Picture_4.jpeg)

The one

### One tile to simulate them all

![](_page_198_Figure_1.jpeg)

Intrinsically universal tile set

Damien Woods

#### http://oscarfarellano.deviantart.com/ 44

The one

### One tile to simulate them all

![](_page_199_Figure_1.jpeg)

Intrinsically universal tile set

**Damien Woods** 

![](_page_199_Picture_4.jpeg)

The one

### Magic dust

![](_page_200_Picture_1.jpeg)

http://nighthawk101stock.deviantart.com/

![](_page_201_Figure_1.jpeg)

Intrinsic universality...requires cooperation. Meunier, Patitz, Summers, Theyssier, Winslow, Woods. SODA 2014

![](_page_202_Figure_1.jpeg)

![](_page_203_Figure_1.jpeg)

### Acknowledgements

![](_page_204_Picture_1.jpeg)

Jack Lutz

Matthew Patitz

Dave Doty Scott Summers

Robert Schweller

![](_page_204_Picture_7.jpeg)

Erik Demaine

![](_page_204_Picture_9.jpeg)

Martin Demaine

![](_page_204_Picture_11.jpeg)

Sándor Fekete

![](_page_204_Picture_13.jpeg)

Pierre-Étienne Meunier

![](_page_204_Picture_15.jpeg)

Guillaume Theyssier

![](_page_204_Picture_17.jpeg)

Andrew Winslow
And many others ....

# Open question: Probabilistically fair intrinsically universal tile set?

![](_page_206_Picture_0.jpeg)

## Structure

- Monday (lecture 1). Some high-level motivations, basic algorithmic selfassembly models (definitions) and very recent results on implementing algorithmic DNA nanotube circuits, a self-assembly model, in the wet-lab
- Tuesday (lecture 2): DNA sequence design and results on the DNA nanotube circuit model
- Wednesday (lecture 3). Complexity theory for self assembly.
  - Theorem: The (cooperative, temperature >= 2) abstract tile assembly model is intrinsic universal
- Thursday (lecture 4). Complexity theory for self assembly.
  - **Theorem**: The noncooperative (temperature 1) abstract tile assembly model does not simulate the cooperative model

![](_page_207_Picture_7.jpeg)

- Temperature 1 tile assembly systems:
  - Tile binds to an assembly if  $\geq$  1 side match
  - Snakes on a plane

![](_page_208_Figure_4.jpeg)

![](_page_208_Picture_5.jpeg)

- Temperature 1 tile assembly systems:
  - Tile binds to an assembly if  $\geq$  1 side match
  - Snakes on a plane

![](_page_209_Figure_4.jpeg)

![](_page_209_Picture_5.jpeg)

- Temperature 1 tile assembly systems:
  - Tile binds to an assembly if  $\geq$  1 side match
  - Snakes on a plane

![](_page_210_Figure_4.jpeg)

- Temperature 1 tile assembly systems:
  - Tile binds to an assembly if  $\geq$  1 side match
  - Snakes on a plane

![](_page_211_Figure_4.jpeg)

- Temperature 1 tile assembly systems:
  - Tile binds to an assembly if  $\geq$  1 side match
  - Snakes on a plane

![](_page_212_Figure_4.jpeg)

• Temperature 1 tile assembly systems:

- Tile binds to an assembly if  $\geq$  1 side match
- Snakes on a plane

![](_page_213_Figure_4.jpeg)

Pic credit: Scott Summers & Dave Doty

![](_page_214_Figure_0.jpeg)

![](_page_215_Figure_0.jpeg)
# Is temperature 1 computationally weak?

Binding on 1 side seems much weaker than binding on 2 sides ... right?

- It has been conjectured (since 2000) that temperature 1 systems are computationally "weak"
  Rothemund, Winfree. STOC 2000
- Some partial negative results:
  - Temperature 1 systems that build fully connected n x n squares require at least n<sup>2</sup> tile types
    Rothemund, Winfree. STOC 2000
  - Pumpable temperature 1 systems produce periodic structures Doty, Patitz, Summers. TCS 2011
  - Temperature 1 with **no mismatches** require 2*n*-1 tile types to assemble an *n* x *n* square

Manuch, Stacho, Stoll. J Comp. Bio. 2010

- Positive results:
  - 3D deterministic temperature 1 simulates Turing machines
  - 2D temperature 1 simulates Turing machine, but with some error
  - 2D temperature 1 can grow large(r than tile set size) structures

Cook, Fu Schweller. SODA 2011 Adleman et al FOCS 2002 Meunier. In submission. 2015

#### On the blackboard: fully-connected square result

**Damien Woods** 

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Manuch, Stacho, Stoll. J Comp. Bio. 2010

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  - 2D temperature 1 can grow large(r than tile set size) structures

Adleman et al FOCS 2002 Meunier. In submission. 2015

- But can temperature 1 aTAM systems simulate cooperative tile assembly?
- Answer: No!

Meunier, Patitz, Summers, Theyssier, Winslow, Woods. SODA 2014

### On the blackboard: fully-connected square result

Damien Woods

## Result

- Theorem 1. There is no tile set U, such that at temperature 1, U simulates all tile assembly systems.
- Theorem 2. There is a 2D temperature 2 tile assembly system T that can not be simulated by any 2D, nor any 3D, temperature 1 tile assembly system.



54

### Temperature 1 can not simulate temperature 2

 We will show that no temperature 1 system simulates the following simple temperature 2 system



# Warm-up: two-seeded system

On the blackboard: a much easier warm-up result

# Simulation definitions

#### Follows:

Definition 3.1 ( $\mathcal{T}$  follows  $\mathcal{U}$ ). We say that  $\mathcal{T}$  follows  $\mathcal{U}$  (under R) if for all  $\alpha', \beta' \in \mathcal{A}[\mathcal{U}]$  where  $\alpha' \to^{\mathcal{U}} \beta'$ , it is the case that  $R^*(\alpha') \to^{\mathcal{T}} R^*(\beta')$ , and  $\alpha' \in \mathcal{A}_{\Box}[\mathcal{U}] \implies R^*(\alpha') \in \mathcal{A}_{\Box}[\mathcal{T}]$ .

#### Models:

Definition 3.2 (nicely fuzzy). We say that  $\mathcal{U}$  is nicely fuzzy with respect to T if for all  $\alpha'' \in \mathcal{A}[\mathcal{U}]$  there exists  $\alpha' \in \mathcal{A}[\mathcal{U}]_{\text{fuzz-free}}$  such that  $\alpha' \to^{\mathcal{U}} \alpha''$ , where  $R^*(\alpha'') = R^*(\alpha') = \alpha \in \mathcal{A}^T$ .

Definition 3.4 ( $\mathcal{U}$  models  $\mathcal{T}$ ). We say that  $\mathcal{U} = (U, \sigma_{\mathcal{U}}, \tau_{\mathcal{U}})$  models  $\mathcal{T} = (T, \sigma_{\mathcal{T}}, \tau_{\mathcal{T}})$  (under R) if:

- (1)  $\mathcal{U}$  is nicely fuzzy with respect to T, and
- (2)  $R^*(\sigma_{\mathcal{U}}) = \sigma_{\mathcal{T}}$ , and
- (3) for all  $\alpha, \beta \in \mathcal{A}[\mathcal{T}]$  such that  $\alpha \to^{\mathcal{T}} \beta$  it is the case that for all  $\alpha' \in \mathcal{A}[\mathcal{U}]_{\text{fuzz-free}}$  where  $R^*(\alpha') = \alpha$ , there exists  $\beta' \in \mathcal{A}[\mathcal{U}]_{\text{fuzz-free}}$  such that  $\alpha' \to^{\mathcal{U}} \beta'$  and  $R^*(\beta') = \beta$ .

#### Simulates:

Definition 3.5. We say that  $\mathcal{U}$  simulates  $\mathcal{T}$  (under R) if  $\mathcal{T}$  follows  $\mathcal{U}$  (under R) and  $\mathcal{U}$  models  $\mathcal{T}$  (under R).

#### Damien Woods

### Temperature 1 is not IU for the aTAM

 First we prove a simple and general pumping lemma for tile assembly at any temperature (the window movie lemma)



 We then use this pumping lemma to "fool" any claimed temperature 1 simulator into exposing its inability to simulate cooperation



d

# Temperature 1 is not IU for the aTAM

- There are simple temperature 2 systems that can not be simulated by any temperature 1 system Meunier, Patitz, Summers, Theyssier, Winslow, Woods. SODA 2014
- First fully-general negative result on temperature 1 (i.e. no restrictions on the model)
- This negative result holds in 2D and 3D

Cook, Fu, Schweller. SODA 2011

- Recall: Deterministic 3D temperature 1 systems can simulate Turing machines!
- So these Turing-universal (powerful!) tile assembly systems can not simulate tile assembly
- Turing universal algorithmic behaviour in self-assembly provably does not imply the ability to simulate arbitrary algorithmic self-assembly processes. Temp 1 3D can compute, but can't handle geometry
- The proof had almost zero "temperature 1 craziness"
  - Ongoing work: with Pierre-Étienne Meunier, & by Damien Regnault and Pierre-Étienne Meunier towards showing other negative results on temperature 1

### Acknowledgements



Pierre-Étienne Meunier



Matthew Patitz



Scott Summers



Andrew Winslow



Guillaume Theyssier

Meunier, Patitz, Summers, Theyssier, Winslow, Woods. SODA 2014

