Tales of the Collatz process in binary (and ternary) (and senary) Tristan Stérin, Damien Woods Hamilton Institute & Computer Science Department Maynooth University, Ireland







Hamilton Institute





ERC No 772766, SFI 18/ERCS/5746

We present two papers

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020

Orthogonal results but same motivation

- I. Motivation: why study the Collatz process in binary?
- II. Results:
 - A. Characterizing the binary representation of ancestors in the Collatz process
 - B. The Collatz process can solve base conversion $3 \rightarrow 2$ (not in AC0)

The Collatz map

- One of the simplest example of *Switched Affine Maps*. Formally introduced in the 60s based on Lothar Collatz's work from the 30s.

J. C. Lagarias. "The 3x+1 problem: an annotated blibliography". 2003, 2006

$$C(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2\\ 3x+1 & \text{if } x \equiv 1 \mod 2\\ & x \in \mathbb{N} \end{cases}$$

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Examples:
$$C(8) = 4$$
 $C(3) = 10$

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Examples:
$$C(8) = 4$$
 $C(3) = 10$
 $C(5) = 16$

The Collatz map, slightly fast forwarded

- One of the simplest example of *Switched Affine Maps*. Formally introduced in the 60s based on Lothar Collatz's work from the 30s.

$$T(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2\\ (3x+1)/2 & \text{if } x \equiv 1 \mod 2\\ & x \in \mathbb{N} \end{cases}$$

$$T(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2\\ (3x+1)/2 & \text{if } x \equiv 1 \mod 2 \end{cases}$$

Let's iterate: 34,

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Let's iterate: 34, 17, 26, 13, 20, 10, 5, 8, 4, 2, 1

$$T(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2\\ (3x+1)/2 & \text{if } x \equiv 1 \mod 2 \end{cases}$$

Let's iterate: 34, 17, 26, 13, 20, 10, 5, 8, 4, 2, **1**, 2, **1**, 2, **1**, ...



The Collatz conjecture $T(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2 \\ (3x+1)/2 & \text{if } x \equiv 1 \mod 2 \end{cases}$ Let's iterate: 34, 17, 26, 13, 20, 10, 5, 8, 4, 2, 1, 2, 1, 2, 1, ...

Collatz conjecture: All strictly positive integers reach **1** under the action of the Collatz process.

The Collatz conjecture $T(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2 \\ (3x+1)/2 & \text{if } x \equiv 1 \mod 2 \end{cases}$ Let's iterate: 34, 17, 26, 13, 20, 10, 5, 8, 4, 2, 1, 2, 1, 2, 1, ...

Collatz conjecture: All strictly positive integers reach **1** under the action of the Collatz process.

Two seemingly independent components:

- 1. Cyclic conjecture: the only strictly positive cycle is 2,1,2,1,...
- 2. Non-divergence conjecture: No Collatz trajectory on the strictly positive integers diverges.

As of 2020, computer tested up to 2^67

D. Barina. The Journal of Supercomputing, 2020.

$$T(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2\\ (3x+1)/2 & \text{if } x \equiv 1 \mod 2\\ x \in \mathbb{N} \end{cases}$$

- $\mathbb{N}(+, \times)$ is a semiring
- There is a "well-defined" parity function $\epsilon : \mathbb{N} \to \mathbb{Z}/2\mathbb{Z}$

 $\epsilon(x \times y) = \epsilon(x) \times \epsilon(y)$

- It is an homomorphism: $\epsilon(1) = 1$ $\epsilon(x + y) = \epsilon(x) + \epsilon(y)$
- It satisfies:

$$\epsilon(x) = 0 \Leftrightarrow 2 \mid x$$

$$T(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2\\ (3x+1)/2 & \text{if } x \equiv 1 \mod 2 \end{cases}$$

$$x\in\mathbb{Z}$$
 Parity function:
 $\epsilon(x)=x ext{ mod } 2$

$$T(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2\\ (3x+1)/2 & \text{if } x \equiv 1 \mod 2 \end{cases}$$

 $x \in \mathbb{Z}$ Parity function: $\epsilon(x) = x \mod 2$

Three known strictly negative cycles:

No known divergent integers

$$T(x) = \begin{cases} x/2 & \text{if } x \equiv 0 \mod 2\\ (3x+1)/2 & \text{if } x \equiv 1 \mod 2\\ x \in \{\frac{a}{b} \mid b \text{ odd }\} \subset \mathbb{Q} \quad \stackrel{\text{Parity function:}}{\epsilon(x) = a \mod 2} \end{cases}$$

Examples:

-2/7, -1/7, 2/7, 1/7, **5/7**, 11/7, 20/7, 10/7, **5/7**, ... 5/3, 3, 5, 8, 4, **2**, 1, **2**, 1, **2**, ...

Lagarias' Periodicity Conjecture: J. C. Lagarias. *The American Mathematical Monthly*, 1985. Every odd-denominator rational reaches a cycle.

$$egin{aligned} x \in \mathbb{N} & x \in \mathbb{Z} & x \in \{rac{a}{b} \mid b ext{ odd }\} \subset \mathbb{Q} \ \epsilon(x) = x ext{ mod } 2 & \epsilon(x) = x ext{ mod } 2 & \epsilon(x) = a ext{$$

$$egin{aligned} x \in \mathbb{N} & x \in \mathbb{Z} & x \in \{rac{a}{b} \mid b ext{ odd }\} \subset \mathbb{Q} \ \epsilon(x) = x ext{ mod } 2 & \epsilon(x) = a ext{ mod } 2 \ \end{aligned}$$
 $\epsilon(x) = 0 \Leftrightarrow 2 \mid x$

Binary decomposition algorithm (i.e., on N, the standard binary encoding)

Binary decomposition algorithm (unbounded)

$$\begin{array}{ccc} x \in \mathbb{N} & x \in \mathbb{Z} & x \in \left\{ \frac{a}{b} \mid b \text{ odd } \right\} \subset \mathbb{Q} \\ \epsilon(x) = x \bmod 2 & \epsilon(x) = x \bmod 2 & \epsilon(x) = a \bmod 2 \end{array}$$

26

while true { Store $\epsilon(x)$ $x := (x - \epsilon(x))/2$ }

$$\begin{array}{ccc} x \in \mathbb{N} & x \in \mathbb{Z} & x \in \left\{ \begin{array}{c} a \\ b \end{array} \mid b \text{ odd} \right\} \subset \mathbb{Q} \\ \epsilon(x) = x \bmod 2 & \epsilon(x) = x \bmod 2 & \epsilon(x) = a \bmod 2 \end{array}$$

...0000011010

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while true {
Store
$$\epsilon(x)$$

 $x := (x - \epsilon(x))/2$







- $\mathbb{Z}_2(+, \times)$ is a ring (uncountable! unordered!)

"2-adic integers" = the set of semi-infinite binary strings

$$\mathbb{Z}_2$$

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$$\mathbb{Z}_2$$

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 - + ...0000001 = ...0000000

$$\mathbb{Z}_2$$

"2-adic integers" = the set of semi-inifinite binary strings

- $\mathbb{Z}_2(+, imes)$ is a ring (uncountable!, unordered!) - $\epsilon(x) = \mathrm{LSB}(x)$ is a suitable parity function $\epsilon(x) = 0 \Leftrightarrow 2 \mid x$



"2-adic integers" = the set of

semi-inifinite binary strings

We can run the Collatz process in \mathbb{Z}_2

Lagarias Periodicity Conjecture:

In \mathbb{Z}_2 , an eventually periodic input to the Collatz process yields to an evenutally periodic Collatz sequence.

J. C. Lagarias, The American Mathematical Monthly, 1985.

No, the Collatz problem is about half-infinite binary strings. (for the scope of this talk)

No, the Collatz problem is about half-infinite binary strings. (for the scope of this talk)

In that context, it is *crucial* to understand the action of the Collatz process in binary.

Computer Science w.r.t the Collatz problem

Three fields (at least) are concerned by the Collatz problem:

Number Theory

Computer Science

Logic
Three fields (at least) are concerned by the Collatz problem:

Number Theory

Computer Science

Logic

- By its formulation
- The approach led to numerous results:

J. C. Lagarias. "The 3x+1 problem: an annotated blibliography". 2003, 2006

- No strictly positive cycle of length < 17.026.679.261 S. Eliahou. Discrete Mathematics, 1993.
- "Almost all strictly positive integers almost lead to 1"
 T. Tao. preprint, 2019.

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Computer Science

Logic

- By its formulation By its formulation in \mathbb{Z}_2
 - Generalised Collatz Maps are Turing complete (exponential slow down): J.H Conway. Number Theory Conference, 1972.

 $G(x) = \begin{cases} x/5+2 & \text{if } x \equiv 0 \mod 5\\ 7(x-1)/5+3 & \text{if } x \equiv 1 \mod 5\\ 2(x-2)/5+1 & \text{if } x \equiv 2 \mod 5\\ 6(x-3)/5+4 & \text{if } x \equiv 3 \mod 5\\ (x-4)/5 & \text{if } x \equiv 4 \mod 5 \end{cases}$ P. Koiran and C. Moore. TCS, 199.

Three fields (at least) are concerned by the Collatz problem:

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Logic

 Peano-independence in discrete dynamical systems. Ex:
 72. Goodstein sequences.

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Peano-independence in discrete dynamical systems. Ex: Goodstein sequences.

Is the Collatz conjecture true?

Computational power of the Collatz process? Is it Turing complete?

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Peano-independence in discrete dynamical systems. Ex: Goodstein sequences.

Is the Collatz conjecture true?

Computational power of the Collatz process?

Is it Turing complete?

Peano or ZFC independence?

Three fields (at least) are concerned by the Collatz problem:

Number Theory

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Computational power of the Collatz process? Is it Turing complete?

Computer Science

- By its formulation in \mathbb{Z}_2
- Generalised Collatz Maps are Turing complete (exponential slow down)



- Computational power of the Collatz process? Is it Turing complete?
- Complexity of prediction problem?



Motivation

- Computational power of the Collatz process? Is it Turing complete?
- Complexity of prediction problem?

Two operations to understand in binary:

- x/2
- 3x+1

Two operations to understand in binary:

x/2 corresponds to shifting the binary representation to the right:

What about 3x + 1 in binary?

3x + 1 = x + (2x + 1)

Meaning: "x + (left_shift(x) + 1)"

Take x = 9 (1001 in binary):

+ 1001

What about 3x + 1 in binary?

3x + 1 = x + (2x + 1)

Meaning: "x + (left_shift(x) + 1)"

Take x = 9 (1001 in binary):

1001 **10010**

What about 3x + 1 in binary?

3x + 1 = x + (2x + 1)

Meaning: "x + (left_shift(x) + 1)"

Take x = 5 (1001 in binary):

1001 **10011**

What about 3x + 1 in binary?

3x + 1 = x + (2x + 1)

Meaning: " $x + (left_shift(x) + 1)$ "

Take x = 9 (1001 in binary):

The Collatz process in binary 3x + 1 = x + (2x + 1) Meaning: 1001 "Each bit of the in

+ 10011 = 11100 "Each bit of the input gets added to its right neighbour and the potential carry on that neighbour."

1001

1001 1001<mark>1</mark> 11100

1001 1001<mark>1</mark> 11100



1001 1001<mark>1</mark> 11100



1001 1001<mark>1</mark> 11100



1001 + 1001<mark>1</mark> = 11100



1001 1001<mark>1</mark> 11100



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..0010010 1100

1001 1001<mark>1</mark> 11100 "Each bit of the input gets added to its right neighbour and the potential carry on that neighbour."

.0010010110010

1001 1001<mark>1</mark> 11100 "Each bit of the input gets added to its right neighbour and the potential carry on that neighbour."

..0010010 ..011000

The Collatz process in binary 3x + 1 = x + (2x + 1)Meaning: "Each bit of the input gets added to its right neighbour 1001 and the potential carry on that neighbour." 11100 ..0010010...011100 3*9 + 1 = 28

Caption read : write



The 3x+1 binary FST














The 3x+1 operation is 2-automatic





The 3x+1 operation is 2-automatic





The 3x+1 operation is 2-automatic





The 3x, 3x+1 and 3x+2 operations are 2-automatic



The 3x, 3x+1 and 3x+2 operations are 2-automatic



"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020 Regular expressions

Base 2 \rightarrow Base 3

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 642:

642

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 642:

321

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 642:



"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 321:

1:0

1:1

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 642:



"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 642:



"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 642:

1010000010 321 1111000100 241 0:10:01011010100 0 181 1:10:11:00:01000100000 544 Ō 1:0

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 642:

1010000010 321 1111000100 241 0:10:01011010100 0 181 1:10:11:00:01000100000 17 ō 1:0

642 is an ancestor of 17

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 642:



642 is an ancestor of 17 at "**odd-distance**" k = 3

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Let's run the Collatz process on 642:



Question:

What is the structure of all ancestors of 17 at "odd-distance" k = 3?

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

United States Let's run the Collatz process on 642: Download PDF Q Find Prior Art ∑ Similar 1010000010 Inventor: Mathieu Ciet, Augustin J. Farrugia, Thomas Icart 321 Current Assignee : Apple Inc 1111000100 Worldwide applications 241 2011 · US 10110101 Application US13/308,452 events ⑦ 181 2011-11-01 • Priority to US201161554411P 2011-11-30 • Application filed by Apple Inc 10001000 17 2011-11-30 Priority to US13/308,452 2011-11-30 Assigned to APPLE INC. @ Question: 2013-05-02 • Publication of US20130108038A1

What is the structure of all ancestors of 17 at "odd-distance" k = 3?

APPLE

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020



Let's denote this set by Pred(k,x)



J. Shallit and D. A. Wilson. Bulletin of the EACTS, 1992.

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Question:

What is the structure of all ancestors of x at "odd-distance" k?

THEOREM:

Pred(k,x) is regular

J. Shallit and D. A. Wilson. Bulletin of the EACTS, 1992.



Doubly exponential regular expressions

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Question:

What is the structure of all ancestors of x at "odd-distance" k?

THEOREM:

Pred(k,x) is regular



"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Example: ancestors of 17 at distance 3.

(101000010010111101) * 10100(000111) * 00(01) * 01(0) *

```
642 <u>101000010</u>
1111000100
1011010100
<u>10001</u>00000
17
```

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Example: ancestors of 17 at distance 3.

(101000010010111101) * 10100(000111) * 00(01) * 01(0) *

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Example: ancestors of 17 at distance 3. There are 12 families.

(101000010010111101)*10100(000111)*00(01)*01(0)*(101000010010111101)*1010000(100011)*100(01)*01(0)*(101000010010111101)*10100001001011110(110001)*1100(01)*01(0)*(101000010010111101)*1010000100101(111000)*11100(01)*01(0)*(101000010010111101)*10100001001(011100)*011100(01)*01(0)*(101000010010111101)*1(001110)*0(01)*01(0)*(101000010010111101)*10100(000111)*0001(10)*1(0)*(101000010010111101)*1010000(100011)*10001(10)*1(0)*(101000010010111101)*10100001001011110(110001)*110001(10)*1(0)*(101000010010111101)*1010000100101(111000)*1(10)*1(0)*(101000010010111101)*10100001001(011100)*01(10)*1(0)*(101000010010111101)*1(001110)*001(10)*1(0)*

Code: https://github.com/tcosmo/coreli

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Result: ancestors of x at distance k. There are $2^{k-1}3^{\frac{(k-1)(k-2)}{2}}$ families.

▶ Theorem 1. For all $x \in \mathbb{N}$, for all $k \in \mathbb{N}$ there exists a regular expression $\operatorname{reg}_k(x)$ that defines $\mathcal{E}\operatorname{Pred}_k(x)$. The regular expression $\operatorname{reg}_k(x)$ is structured as a tree with $2^k 3^{k(k-1)/2}$ branches, alphabetic width $O(2^k 3^{k(k+1)/2})$ and star height equal to 1.

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

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▶ **Theorem 1.** For all $x \in \mathbb{N}$, for all $k \in \mathbb{N}$ there exists a regular expression $\operatorname{reg}_{k}(x)$ that defines $\mathcal{E}\operatorname{Pred}_{k}(x)$. The regular expression $\operatorname{reg}_{k}(x)$ is structured as a tree with $2^{k}3^{k(k-1)/2}$ branches, alphabetic width $O(2^{k}3^{k(k+1)/2})$ and star height equal to 1. Erratum: k := k-1

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Example: ancestors of <u>17</u> at distance 3. There are **12 families.**

 Complexity: at least as hard as solving discrete logarithm in Z/3^kZ. I.e: find `i` such that x = 2⁻i mod [3^k].

Take: **len**(**)**-2, here 5-2 = 3

 $2^{-}(3) = 14^{3} = 2744 = 17 \mod 3^{3}$

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Example: ancestors of <u>17</u> at distance 3. There are **12 families.**

- Complexity: at least as hard as solving discrete logarithm in Z/3^kZ. I.e: find `i` such that x = 2⁻i mod [3^k].
- `Seeds` are mysterious:

Equivalent definitions:

- \circ Repetend of 1/3^k in \mathbb{Z}_2
- \circ Parity bits of 1/2^k in $\mathbb{Z}/3^k\mathbb{Z}$

$$2 \cdot 3^{k-1}$$

• Binary expansion of $\frac{2}{3^k}$

Seeds approach Full Complexity:

J. Lòpez and P. Stoll. Integers, 2012

"Binary representation of ancestors in the Collatz graph", T. Stérin, RP 2020

Example: ancestors of <u>17</u> at distance 3. There are **12 families.**

Futur work:

Efficiently finding the smallest ancestor at distance k?

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020

Finite state automaton



"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020

2D Cellular automaton



"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020

2D Cellular automaton



I. Korec. Mathematica Slovaca, 1992.

Quasi base 2 CA

T. Cloney, E. Goles and G. Vichiniac. Complex Systems, 1987.

Eric Goles, Nicolas Ollinger, and Guillaume Theyssier.













"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020



Column by column

0:1

 $\overline{1}:\overline{1}$

^{`1120`} **= 4**2
"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020



Column by column

0:1

 $\overline{1}:\overline{1}$

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020



 $\overline{1}:\overline{1}$

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020





(c) Horizontal applications of the local CQCA rule simulate the 3x + 1 FST.



(a) The 3x + 1 binary FST.



(d) Vertical applications of the local rule simulate the x/2 FST.



(b) The x/2 base 3' FST.

Dual Transducers

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020





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0:0



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Dual Transducers

Computational power of the Collatz process? Is it Turing complete?

Complexity of prediction problem?

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020

Motivation

Complexity of prediction:



ACO: polysize, constant depth, arbitrary fanin, Bolean circuits NC1: polysize, log depth, fanin <= 2, Bolean circuits

AC0 != P, problem PARITY not in AC0

Theorem 28 (arxiv numbering):

The upper z-rectangle prediction problem is in NC1 and not in AC0.

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020

Complexity of prediction:



process? Is it Turing complete? Complexity of prediction problem?



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- Computational power of the Collatz process? Is it Turing complete?
- Complexity of prediction problem?

NC1: polysize, log depth, fanin <= 2, Bolean circuits (c) (d) AC0 != P, problem PARITY not in AC0 upper *z*-rectangle 0 Б base 3-to-2 base Theorem 28 (arxiv numbering): $T^n(z)$ The upper z-rectangle prediction problem co is in <u>NC1</u> and not in <u>AC0</u>. z in base 2; n bits apply Tⁿ to z (z) in base W. Hesse, E. Allender and D. Barrington. Π Journal of Computer Science and System Science, 2002. 0 0 0 TOO lower *z*-rectangle Answering if a base 3-encoded number

ACO: polysize, constant depth, arbitrary fanin, Bolean circuits

is odd or even reduces to PARITY.

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020

Future Work

Motivation

 Computational power of the Collatz process? Is it Turing complete?
Complexity of prediction problem?



Theorem 28 (arxiv numbering):

The upper z-rectangle prediction problem is in $\underline{NC1}$ and not in $\underline{AC0}$.

What about the lower z-rectangle? In P, not in AC0. Can we get matching bounds?

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020

Motivation Computational power of the Collatz process? Is it Turing complete? Future Work - Complexity of prediction problem?

What about this diagonal?

Magenta column: x = `1120102` = 1145 Blue line: x = `10001111001` = 1145

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Future Work

Motivation

 Computational power of the Collatz process? Is it Turing complete?
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What about this diagonal?

Magenta column: x = `1120102` = 1145 Blue line: x = `10001111001` = 1145

Diagonal is x in base 6

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Future Work

Motivation

 Computational power of the Collatz process? Is it Turing complete?
Complexity of prediction problem?



Need to further understand this encoding

- Link with Chinese Remainder
- Link with $Z2 \times Z3 = Z6$

$$x = 5145 = 1145 \begin{array}{c|cccc} 0 & 0 & \overline{0} & 1 & \overline{1} & \overline{1} \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 2 & 3 & 4 & 5 & 119 \end{array}$$

"The Collatz process embeds a base conversion algorithm", T. Stérin and D. Woods, RP 2020



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Conclusion

- Using base 2/3/6 Finite State Automaton and Cellular Automaton descriptions of the Collatz process brings it in the scope of Computer Science.
- In that context we have a sound framework to ask:
 - What is the computational power of the Collatz process? Is it Turing complete?
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Conclusion

- Using base 2/3/6 Finite State Automaton and Cellular Automaton descriptions of the Collatz process brings it in the scope of Computer Science.
- In that context we have a sound framework to ask:
 - What is the computational power of the Collatz process? Is it Turing complete?
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So, what's next?

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